# MATH 567: Mathematical Techniques in Data Science Clustering II

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This lecture is based on U. von Luxburg, A Tutorial on Spectral Clustering, Statistics and Computing, 17 (4), 2007.

Overview of spectral clustering:

• Construct a *similarity matrix* measuring the similarity of pairs of objects (e.g.  $s_{ij} = \exp(-||x_i - x_j||^2/(2\sigma^2)))$ ).

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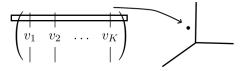
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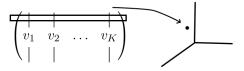


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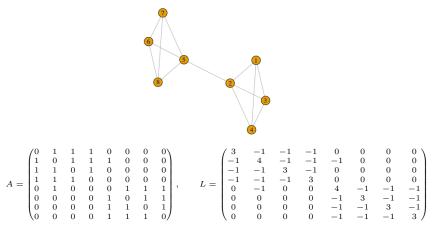
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#### Example

Let us try to cluster the following graph:



#### We have:

 $v_2 = (-0.3825277, -0.2470177, -0.3825277, -0.3825277, 0.2470177, 0.3825277, 0.3825277, 0.3825277)^T.$ 

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**3** 0 is an eigenvalue of L with associated constant eigenvector 1. **Proof:** To prove (1),

$$f^{T}Lf = f^{T}Df - f^{T}Wf = \sum_{i=1}^{n} d_{i}f_{i}^{2} - \sum_{i,j=1}^{n} w_{ij}f_{i}f_{j}$$
$$= \frac{1}{2} \left( \sum_{i=1}^{n} d_{i}f_{i}^{2} - 2\sum_{i,j=1}^{n} w_{ij}f_{i}f_{j} + \sum_{j=1}^{n} d_{j}f_{j}^{2} \right)$$
$$= \frac{1}{2} \sum_{i,j=1}^{n} w_{ij}(f_{i} - f_{j})^{2}.$$

(2) follows from (1). (3) is easy.

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**Proof:** If f is an eigenvector associated to  $\lambda = 0$ , then

$$0 = f^T L f = \sum_{i,j=1}^n w_{ij} (f_i - f_j)^2.$$

It follows that  $f_i = f_j$  whenever  $w_{ij} > 0$ . Thus f is constant on the connected components of G. We conclude that the eigenspace of 0 is contained in  $\operatorname{span}(\mathbb{1}_{A_1}, \ldots, \mathbb{1}_{A_k})$ . Conversely, it is not hard to see that each  $\mathbb{1}_{A_i}$  is an eigenvector associated to 0 (write L in block diagonal form). **Proposition:** The normalized Laplacians satisfy the following properties:

**1** For every  $f \in \mathbb{R}^n$ , we have

$$f^T L_{\text{sym}} f = \frac{1}{2} \sum_{i,j=1}^n w_{ij} \left( \frac{f_i}{\sqrt{d_i}} - \frac{f_j}{\sqrt{d_j}} \right)^2.$$

- 2  $\lambda$  is an eigenvalue of  $L_{\rm rw}$  with eigenvector u if and only if  $\lambda$  is an eigenvalue of  $L_{\rm sym}$  with eigenvector  $w = D^{1/2}u$ .
- $\lambda$  is an eigenvalue of  $L_{rw}$  with eigenvector u if and only if  $\lambda$  and u solve the generalized eigenproblem  $Lu = \lambda Du$ .

**Proof:** The proof of (1) is similar to the proof of the analogous result for the unnormalized Laplacian. (2) and (3) follow easily by using appropriate rescalings.

**Proposition:** Let G be an undirected graph with non-negative weights. Then:

- The multiplicity k of the eigenvalue 0 of both L<sub>sym</sub> and L<sub>rw</sub> equals the number of connected components A<sub>1</sub>,..., A<sub>k</sub> in the graph.
- 2 For  $L_{\rm rw}$ , the eigenspace of eigenvalue 0 is spanned by the indicator vectors  $\mathbb{1}_{A_i}$ ,  $i = 1, \ldots, k$ .
- For  $L_{\rm sym}$ , the eigenspace of eigenvalue 0 is spanned by the vectors  $D^{1/2} \mathbbm{1}_{A_i}$ ,  $i = 1, \ldots, k$ .

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### Graph cuts

- G graph with (weighted) adjacency matrix  $W = (w_{ij})$ .
- We define:

$$W(A,B) := \sum_{i \in A, j \in B} w_{ij}.$$

- |A| := number of vertices in A.
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Given a partition  $A_1,\ldots,A_k$  of the vertices of G, we let

$$\operatorname{cut}(A_1,\ldots,A_k) := \frac{1}{2} \sum_{i=1}^k W(A_i,\overline{A}_i).$$

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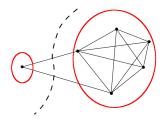
The min-cut problem consists of solving:

$$\min_{\substack{V=A_1\cup\cdots\cup A_k\\A_i\cap A_j=\emptyset\;\forall i\neq j}} \operatorname{cut}(A_1,\ldots,A_k).$$

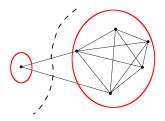
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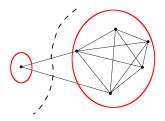


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- We would like clusters to have a reasonably large number of points.
- We therefore modify the min-cut problem to enforce such constraints.

The two most common objective functions that are used as a replacement to the min-cut objective are:

Q RatioCut (Hagen and Kahng, 1992):

RatioCut
$$(A_1, \ldots, A_k) := \frac{1}{2} \sum_{i=1}^k \frac{W(A_i, \overline{A}_i)}{|A_i|} = \sum_{i=1}^k \frac{\operatorname{cut}(A_i, \overline{A}_i)}{|A_i|}.$$

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In Normalized cut (Shi and Malik, 2000):

$$\operatorname{Ncut}(A_1,\ldots,A_k) := \frac{1}{2} \sum_{i=1}^k \frac{W(A_i,\overline{A}_i)}{\operatorname{vol}(A_i)} = \sum_{i=1}^k \frac{\operatorname{cut}(A_i,\overline{A}_i)}{\operatorname{vol}(A_i)}.$$

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- Note: both objective functions take larger values when the clusters A<sub>i</sub> are "small".
- Resulting clusters are more "balanced".
- However, the resulting problems are NP hard see Wagner and Wagner (1993).

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Given  $A \subset V$ , let  $f \in \mathbb{R}^n$  be given by

$$f_i := \begin{cases} \sqrt{|\overline{A}|/|A|} & \text{if } v_i \in A\\ -\sqrt{|A|/|\overline{A}|} & \text{if } v_i \notin A. \end{cases}$$

#### Relaxing RatioCut

Let L = D - W be the (unnormalized) Laplacian of G. Then

$$\begin{split} f^T L f &= \frac{1}{2} \sum_{i,j=1}^n w_{ij} (f_i - f_j)^2 \\ &= \frac{1}{2} \sum_{i \in A, j \in \overline{A}} w_{ij} \left( \sqrt{\frac{|\overline{A}|}{|A|}} + \sqrt{\frac{|A|}{|\overline{A}|}} \right)^2 + \frac{1}{2} \sum_{i \in \overline{A}, j \in A} w_{ij} \left( -\sqrt{\frac{|\overline{A}|}{|A|}} - \sqrt{\frac{|A|}{|\overline{A}|}} \right)^2 \\ &= W(A, \overline{A}) \left( 2 + \frac{|\overline{A}|}{|A|} + \frac{|A|}{|\overline{A}|} \right) \\ &= W(A, \overline{A}) \left( \frac{|A| + |\overline{A}|}{|A|} + \frac{|A| + |\overline{A}|}{|\overline{A}|} \right) \\ &= |V| \cdot \frac{1}{2} \left( \frac{W(A, \overline{A})}{|A|} + \frac{W(\overline{A}, A)}{|\overline{A}|} \right) \\ &= |V| \cdot \operatorname{RatioCut}(A, \overline{A}). \\ &\text{since } |A| + |\overline{A}| = |V|, \text{ and } W(A, \overline{A}) = W(\overline{A}, A). \end{split}$$

• We showed:

$$f^T L f = \frac{1}{2} \sum_{i,j=1}^n w_{ij} (f_i - f_j)^2 = |V| \cdot \operatorname{RatioCut}(A, \overline{A}).$$

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• Moreover, note that

$$\sum_{i=1}^n f_i = \sum_{i \in A} \sqrt{\frac{|\overline{A}|}{|A|}} - \sum_{i \in \overline{A}} \sqrt{\frac{|A|}{|\overline{A}|}} = |A| \cdot \sqrt{\frac{|\overline{A}|}{|A|}} - |\overline{A}| \cdot \sqrt{\frac{|A|}{|\overline{A}|}} = 0.$$

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• Finally,

$$||f||_2^2 = \sum_{i=1}^n f_i^2 = |A| \cdot \frac{|\overline{A}|}{|A|} + |\overline{A}| \cdot \frac{|A|}{|\overline{A}|} = |V| = n.$$

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Thus, we have showed that the Ratio-Cut problem is equivalent to  $\min_{A \subset V} f^T L f$ subject to  $f \perp 1, ||f|| = \sqrt{n}, f_i$  defined as above.

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The natural relaxation of the problem is to remove the discreteness condition on f and solve

$$\min_{f \in \mathbb{R}^n} f^T L f$$
  
subject to  $f \perp 1, ||f|| = \sqrt{n}$ .

• Using properties of the Rayleigh quotient, it is not hard to show that the solution of

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• How do we get the clusters from  $\tilde{f}$ ? • We could set

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• More generally, we *cluster* the coordinates of f using K-means.

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$$\begin{cases} v_i \in A & \text{ if } f_i \ge 0\\ v_i \in \overline{A} & \text{ if } f_i < 0. \end{cases}$$

• More generally, we *cluster* the coordinates of f using K-means. This is the **unnormalized spectral clustering algorithm** for k = 2.

• Using properties of the Rayleigh quotient, it is not hard to show that the solution of

$$\min_{f \in \mathbb{R}^n} f^T L f$$
  
subject to  $f \perp 1, ||f|| = \sqrt{n}.$ 

is an eigenvector of  $\boldsymbol{L}$  corresponding to the second eigenvalue.

 $\bullet$  Clearly, if  $\tilde{f}$  is the solution of the problem, then

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- More generally, we *cluster* the coordinates of f using K-means. This is the **unnormalized spectral clustering algorithm** for k = 2.
- $\bullet$  The above process can be generalized to  $k\geq 2$  clusters.

#### The unnormalized spectral clustering algorithm:

Unnormalized spectral clustering

Input: Similarity matrix  $S \in \mathbb{R}^{n \times n}$ , number k of clusters to construct.

- Construct a similarity graph by one of the ways described in Section 2. Let W be its weighted adjacency matrix.
- Compute the unnormalized Laplacian L.
- Compute the first k eigenvectors  $u_1, \ldots, u_k$  of L.
- Let  $U \in \mathbb{R}^{n imes k}$  be the matrix containing the vectors  $u_1, \dots, u_k$  as columns.
- For  $i=1,\ldots,n$ , let  $y_i\in\mathbb{R}^k$  be the vector corresponding to the *i*-th row of U.
- Cluster the points  $(y_i)_{i=1,\ldots,n}$  in  $\mathbb{R}^k$  with the  $k\text{-means algorithm into clusters }C_1,\ldots,C_k.$

Output: Clusters  $A_1, \ldots, A_k$  with  $A_i = \{j | y_j \in C_i\}$ .

Source: von Luxburg, 2007.

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Normalized spectral clustering according to Shi and Malik (2000)

Input: Similarity matrix  $S \in \mathbb{R}^{n \times n}$ , number k of clusters to construct.

- Construct a similarity graph by one of the ways described in Section 2. Let  ${\cal W}$  be its weighted adjacency matrix.
- Compute the unnormalized Laplacian L.
- Compute the first k generalized eigenvectors  $u_1, \ldots, u_k$  of the generalized eigenproblem  $Lu = \lambda Du$ .
- Let  $U \in \mathbb{R}^{n \times k}$  be the matrix containing the vectors  $u_1, \ldots, u_k$  as columns.
- For  $i = 1, \ldots, n$ , let  $y_i \in \mathbb{R}^k$  be the vector corresponding to the *i*-th row of U.
- Cluster the points  $(y_i)_{i=1,\ldots,n}$  in  $\mathbb{R}^k$  with the  $k\text{-means algorithm into clusters }C_1,\ldots,C_k$  .

Output: Clusters  $A_1, \ldots, A_k$  with  $A_i = \{j | y_j \in C_i\}$ .

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Source: von Luxburg, 2007.

• Note: The solutions of  $Lu = \lambda Du$  are the eigenvectors of  $L_{\rm rw}$ .

See von Luxburg (2007) for details.

# The normalized clustering algorithm of Ng et al.

• Another popular variant of the spectral clustering algorithm was provided by Ng, Jordan, and Weiss (2002).

• The algorithm uses  $L_{\rm sym}$  instead of L (unnormalized clustering) or  $L_{\rm rw}$  (Shi and Malik's normalized clustering).

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• The algorithm uses  $L_{\rm sym}$  instead of L (unnormalized clustering) or  $L_{\rm rw}$  (Shi and Malik's normalized clustering).

Normalized spectral clustering according to Ng, Jordan, and Weiss (2002)

Input: Similarity matrix  $S \in \mathbb{R}^{n \times n}$ , number k of clusters to construct.

- $\bullet$  Construct a similarity graph by one of the ways described in Section 2. Let W be its weighted adjacency matrix.
- Compute the normalized Laplacian  $L_{\rm sym}$ .
- Compute the first k eigenvectors  $u_1, \ldots, u_k$  of  $L_{\text{sym}}$ .
- Let  $U \in \mathbb{R}^{n \times k}$  be the matrix containing the vectors  $u_1, \ldots, u_k$  as columns.
- Form the matrix T ∈ ℝ<sup>n×k</sup> from U by normalizing the rows to norm 1, that is set t<sub>ij</sub> = u<sub>ij</sub>/(∑<sub>k</sub> u<sup>2</sup><sub>ik</sub>)<sup>1/2</sup>.
- For  $i=1,\ldots,n$ , let  $y_i\in\mathbb{R}^k$  be the vector corresponding to the i-th row of T.
- Cluster the points  $(y_i)_{i=1,\dots,n}$  with the k-means algorithm into clusters  $C_1,\dots,C_k$ . Output: Clusters  $A_1,\dots,A_k$  with  $A_i=\{j|\ y_j\in C_i\}$ .

Source: von Luxburg, 2007.

#### See von Luxburg (2007) for details.