MATH 829: Introduction to Data Mining and Analysis Independent component analysis

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Motivation

 Blind signal separation: separation of a mixture of source signals, without (or with very little) information about the sources and the mixing process.

• Example (the cocktail party problem): isolate a single conversation in a noisy room with many people talking.



Mathematical formulation



$$x_1(t) = a_{11}s_1(t) + a_{12}s_2(t)$$

$$x_2(t) = a_{21}s_1(t) + a_{22}s_2(t)$$

- We have x(t) = As(t), t = 1, ..., T.
- We observe x(t).
- . We don't know what A is (mixing matrix).
- We don't observe s(t).

We want to recover s(t) (and/or A).

- Current formulation is ill-posed: there are multiple ways of mixing signals to get the output.
- We will seek a solution where the components of s are as independent as possible.

Assumptions

Note: Signals can only be recovered up to

- Permutations: we can always permute the s_i's and the row/columns of A to obtain new solutions.
- Scaling: we can always rescale the s_i's and rescale the coefficients in A.

(Not a big deal in most applications.) Other problems?

Problem with Gaussian data:

- Suppose s ∼ N(0_{2×1}, I_{2×2}) (independent Gaussian sources).
- Let x = As where $A \in \mathbb{R}^{2 \times 2}$.
- Then $x \sim N(\mathbf{0}_{2 \times 1}, AA^T)$.
- Let U be an orthogonal matrix, i.e., $UU^T = U^T U = I$.
- Let A' = AU.
- Then $x' = A's \sim N(\mathbf{0}_{2\times 1}, A'A^T) = N(\mathbf{0}_{2\times 1}, AUU^TA^T) = N(\mathbf{0}_{2\times 1}, AA^T).$

Thus, there is no way to statistically differentiate if x was obtained from the mixing matrix A or A'.

We will therefore assume the sources are not Gaussian.

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Independence of the sources

- We seek sources that are as independent as possible.
- Multiple ways to measure independence. For example:
 - O Minimization of mutual information.
- Maximization of non-Gaussianity measures (negentropy, kurtosis, etc.).

Motivation for (2) comes from the central limit theorem: a sum of independent random variables should be "more Gaussian".



To explain the above notions, we briefly discuss some concepts from *information theory*.

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Entropy of a random variable

- Let X be a random variable taking values $x_1, ..., x_N$ with probabilities $P(X = x_i) = p_i$.
- The entropy of X is given by:

$$H(X) = E(-\log p) = -\sum_{i=1}^{N} p_i \log p_i$$

(usually, we take the log in base 2).

• Similarly, if X is a continuous random variable with density f(x), we define:

$$H(X) = -\int f(x) \log f(x) dx$$

The entropy is a measure of the uncertainty or complexity of a random variable.

Example: If X is a (discrete) uniform on $\{1, ..., N\}$, then

$$H(X) = -\sum_{i=1}^{N} \frac{1}{N} \log\left(\frac{1}{N}\right) = \log N$$

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Entropy (cont.)

Example: $X \sim \text{Bernoulli}(p)$, i.e., P(X = 1) = p, P(X = 0) = 1 - p. The more "uncertain" the outcome is, the larger the entropy.



Entropy and information

We would like to define a measure of information I(p) of an event occurring with probability p. This functions should satisfy:

- I(p) ≥ 0.
- I(1) = 0 (the information gained from observing a certain event is 0).
- I(p1p2) = I(p1) + I(p2) (information gained from observing two independent event is sum of information).
- I should be continuous and monotonic.

The above properties imply $I(p) = \log_b \frac{1}{p}$ for some base b. The entropy of X is the **average information** "contained" in X:

$$H(X) = \sum_{i=1}^{N} I(p_i)p_i$$

Entropy and communication

- Suppose we can only transmit 0s and 1s.
- . We need to encode our message (e.g. choose a code for each letter).
- How efficiently can we encore the message?



Example: Our source sends the letters A, B, C, D. Each letter is equally likely to be transmitted.

 $A \rightarrow 00$ $C \rightarrow 10$ W $B \rightarrow 01$ $D \rightarrow 11$ ac

We send on average (actually, exactly!) 2 bits per symbol.

 If the symbols an not equally likely, it is not hard to see that one can do better (i.e., send less bits per symbol on average).

• The entropy provides a lower bound on the average number of bits required per symbol.

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Kullback-Leibler divergence

Given two (discrete) probability distributions P and Q, we define the Kullback-Leibler divergence by

$$D_{KL}(P||Q) := \sum_{i} P(i) \log \frac{P(i)}{Q(i)}$$

Similarly, when P and Q are continuous with densities $p(\boldsymbol{x})$ and $q(\boldsymbol{x})$ respectively, we define

$$D_{KL}(P||Q) := \int p(x) \log \frac{p(x)}{q(x)} dx$$

Intuitive interpretation:

- A source send symbols with distribution P.
- We encode the messages as if the source had distribution Q.
- D_{KL}(P||Q) is the number of supplementary bits per symbol that we send for not using the "right" distribution.

The KL divergence is used as a measure of distance between distributions (note however that $D_{\mathrm{KL}}\left(P||Q\right) \neq D_{\mathrm{KL}}(Q||P)$ in general).

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Mutual information

- (X_1, \ldots, X_n) random vector with distribution $p(x_1, \ldots, x_n)$.
- Let p(x₁),..., p(x_n) denote the marginals of p (i.e., the distribution of each variable X_i).
- Let (Y₁,...,Y_n) have distribution p(x₁)p(x₂)...p(x_n) (so Y_i has the same distribution as X_i, but the Y_is are independent).

The mutual information of (X_1, \ldots, X_n) is given by

$$I(X_1, ..., X_n) = D_{KL}(p(x_1, ..., x_n)||p(x_1)...p(x_n)).$$

- We have I(X,Y) = 0 if and only if X,Y are independent.
- Therefore, $I(X_1, \ldots, X_n)$ provides a numerical measure of how independent random variables are.

Measures of non-Gaussianity

• The kurtosis (from greek κυρτός, "curved") of a random variable with mean $\mu = E(X)$ is given by

$$\operatorname{Kurt}(X) := \frac{E[(X - \mu)^4]}{(E[(X - \mu)^2])^2}.$$

- Measures the "propensity to produce outliers".
- The Gaussian distribution has kurtosis equal to 3.
- Can thus use the "excess kurtosis" Kurt(X) 3 to test for "non-Gaussianity".
- The negent ropy of a random variable X is given by

 $J(X) := H(X_{gauss}) - H(X),$

where X_{gauss} is a Gaussian random variable with the same mean and variance as X.

- Motivated by the fact that the Gaussian distribution has the largest entropy among all continuous distributions with a given mean and variance.
- Therefore, a variable that is "far from a Gaussian" should have a larger negentropy.

The FastICA algorithm

- FastICA (Hyvärinen, 1999) is an efficient and popular algorithm for computing independent components.
- Finds linear combinations maximizing an approximation of the negentropy.
- . The negentropy is replaced by the approximation

$$J(X) \approx [E(G(X)) - E(G(X_{gauss}))]^2$$
,

where $G(x) = \log \cosh(x)$.

Whitening the data

Before the FastICA algorithm is applied, the data needs to be prewhitened.

- Let $X \in \mathbb{R}^{N \times M}$ be the data matrix.
- First center the rows of X:

$$x_{ij} \leftarrow x_{ij} - \frac{1}{M} \sum_{k} x_{ik}.$$

• Next, we want the linearly transform the rows of X so that they become uncorrelated. We seek a linear transformation $L: \mathbb{R}^{N \times M} \rightarrow \mathbb{R}^{N \times M}$ such that

$$\frac{1}{M}L(x)L(x)^{T} = I_{N \times N}.$$

This is easily achieved using the eigendecomposition of the covariance matrix of the centered data X:

$$\frac{1}{M}XX^{T} = UDU^{T}.$$

. Define the whitened data matrix by

$$X_{mbite} := UD^{-1/2}U^T X.$$

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Python - example

We mix two sound files, and recover them using ICA.

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We want to extract independent components of the form w^TX where $w \in \mathbb{R}^N.$

The FastICA algorithm:

The FastICA algorithm

- Find a first direction w₁ maximizing the (approximation of) the negentropy (can use a fixed point method).
- Estimate a second direction w₂ ± w₁ maximizing the (approximation of) the negentropy.

etc..

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